

OPENFABRICS  
ALLIANCE

12<sup>th</sup> ANNUAL WORKSHOP 2016

# KERNEL VERBS API UPDATE

Leon Romanovsky

Mellanox Technologies

[ April 4<sup>th</sup>, 2016 ]



# AGENDA

- Memory registration API
- CQ polling API
- Draining QP
- Generic RDMA WRITE/READ API



# MEMORY REGISTRATION API

# MULTIPLE MEMORY REGISTRATION METHODS

## ■ Physical memory regions (MR)

- Synchronous interface
- Every registration causes to new MR

## ■ Fast memory registration (FMR)

- Fast synchronous interface
- Weak deregistration semantics
- Not widely adopted

## ■ Memory windows (MW)

- Fast, asynchronous interface
- Binds continuous apertures to existing memory regions
  - Not relevant to kernel ULPs
- Removed from kernel API

## ■ Fast memory registration mode (FRWR)

- Asynchronous interface
- Maps blocks of physical memory
- Widely adopted

# MULTIPLE MEMORY REGISTRATION METHODS

## ■ Physical memory regions (MR)

- Synchronous interface
- Every registration causes to new MR



## ■ Fast memory registration (FMR)

- Fast synchronous interface
- Weak deregistration semantics
- Not widely adopted



## ■ Memory windows (MW)

- Fast, asynchronous interface
- Binds continuous apertures to existing memory regions
  - Not relevant to kernel ULPs
- Removed from kernel API



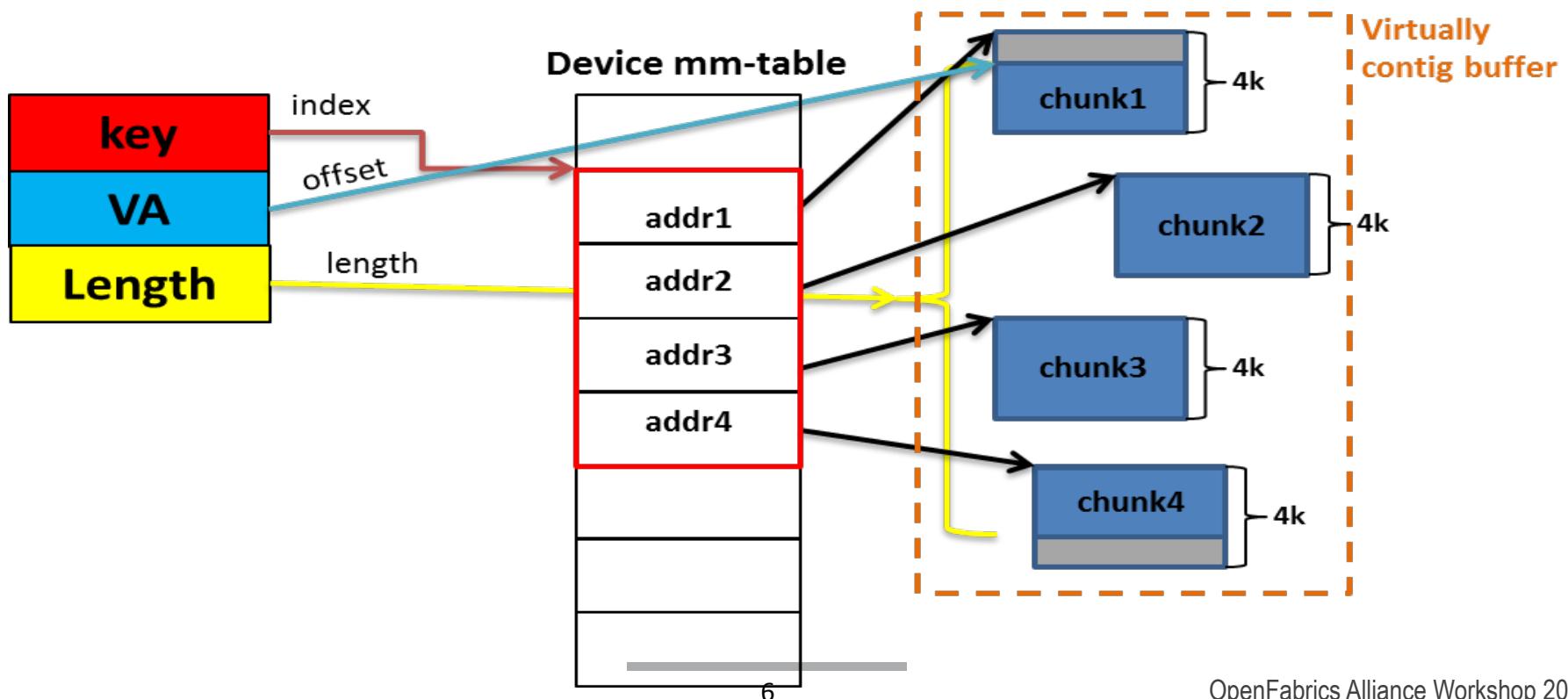
## ■ Fast memory registration mode (FRWR)

- Asynchronous interface
- Maps blocks of physical memory
- Widely adopted



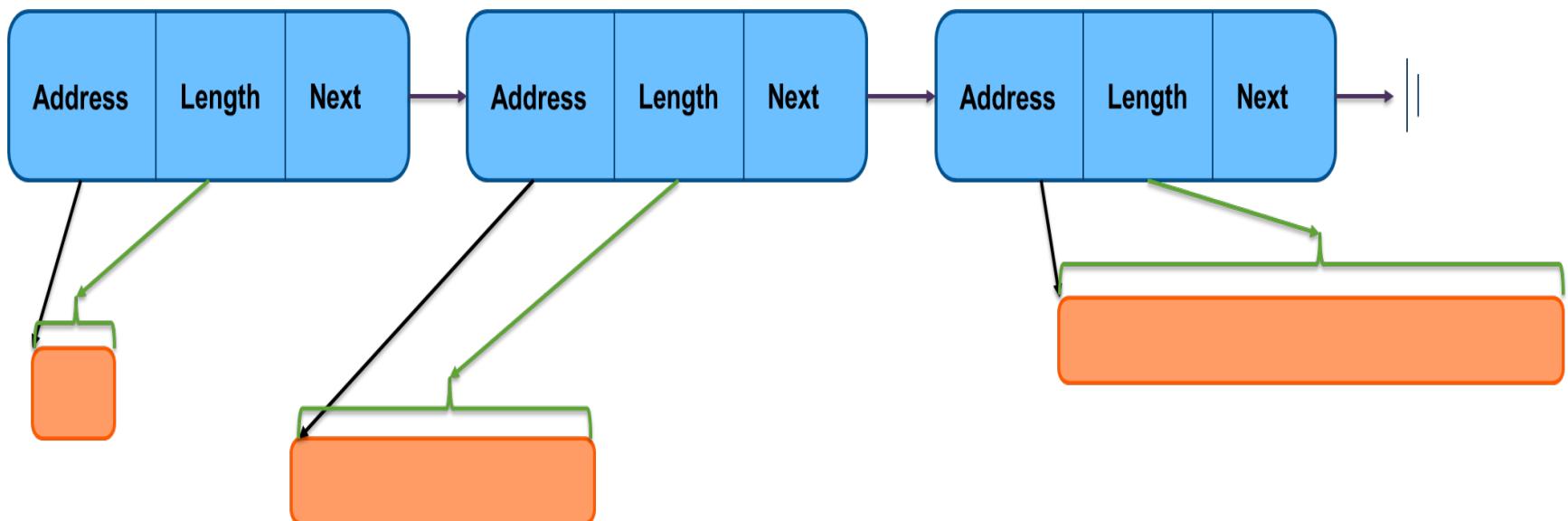
# FRWR MAPPINGS

- All buffers must be block aligned (page shift)
- The first buffer can have first byte offset (FBO)
- The last buffer can end before the end of block (EOB)



# FRWR DRAWBACKS

- Allocation of free memory region and a `fast_reg_page_list`
- Translation from S/G list to page vector for each ULP
- No ability to support arbitrary S/G
  - Each ULP bridged this semantic gap



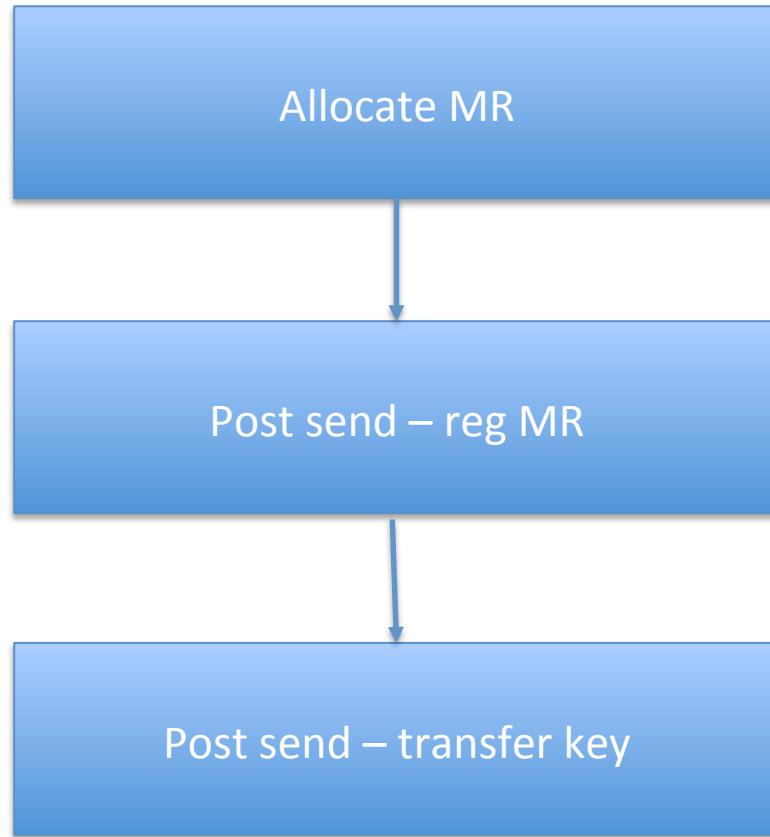
# OLD FRWR API

OP	API
<b>Allocate</b>	<pre>frpl = ib_alloc_fast_reg_page_list() mr = ib_alloc_fast_reg_mr()</pre>
<b>Free</b>	<pre>ib_free_fast_reg_page_list(frpl) ib_dereg_mr(mr)</pre>
<b>Register interface (post_send)</b>	<pre>struct {     u64                                iova_start;     struct ib_fast_reg_page_list *page_list;     unsigned int                      page_shift;     unsigned int                      page_list_len;     u32                                length;     int                                 access_flags;     u32                                rkey; } fast_reg;</pre>
<b>Opcode</b>	IB_WR_FAST_REG_MR

# NEW FRWR API

OP	API
<b>Allocate</b>	mr = ib_alloc_mr(pd, max_num_sg, mr_type)
<b>Free</b>	ib_dereg_mr(mr)
<b>S/G list mapping</b>	nents = ib_map_mr_sg(mr, sg, sg_nents, page_size);
<b>Register interface (post_send)</b>	struct { struct ib_send_wr wr; struct ib_mr *mr; int access; u32 key; } ib_reg_mr;
<b>Opcode</b>	IB_WR_REG_MR

# ULP REGISTRATION FLOW





# CQ POLLING API

# CQ POLLING CONSIDERATIONS

- **Different completion contexts**
  - Kernel threads
  - Work queues
  - Software IRQs
  - Hardware IRQs
- **Fairness between multiple CQs**
- **CQ re-arm policy**
- **Handling missed events**

# CQ POLLING USAGE UNCERTAINTY

- Work requests (WR) context return path
- Work request IDs (wr\_id) (un)reliable
- Multiple sources of post\_send completions
- Polling in batches or one-by-one
- Multiple CPUs (affinity)
- CPU/NUMA locality

# CQ POLLING API

OP	API
<b>Decide on polling type</b>	<pre>enum ib_poll_context {     IB_POLL_DIRECT, /* caller context, no hw completions */     IB_POLL_SOFTIRQ, /* poll from softirq context */     IB_POLL_WORKQUEUE, /* poll from workqueue */ };</pre>
<b>Allocate</b>	<pre>struct ib_cq *ib_alloc_cq(dev, private, nr_cqe, comp_vector, poll_ctx);</pre>
<b>Free</b>	<pre>void ib_free_cq(cq)</pre>
<b>Completion return (filled in WR)</b>	<pre>struct ib_cqe {     void (*done)(struct ib_cq *cq, struct ib_wc *wc); };</pre>

# CQ POLLING SUMMARY

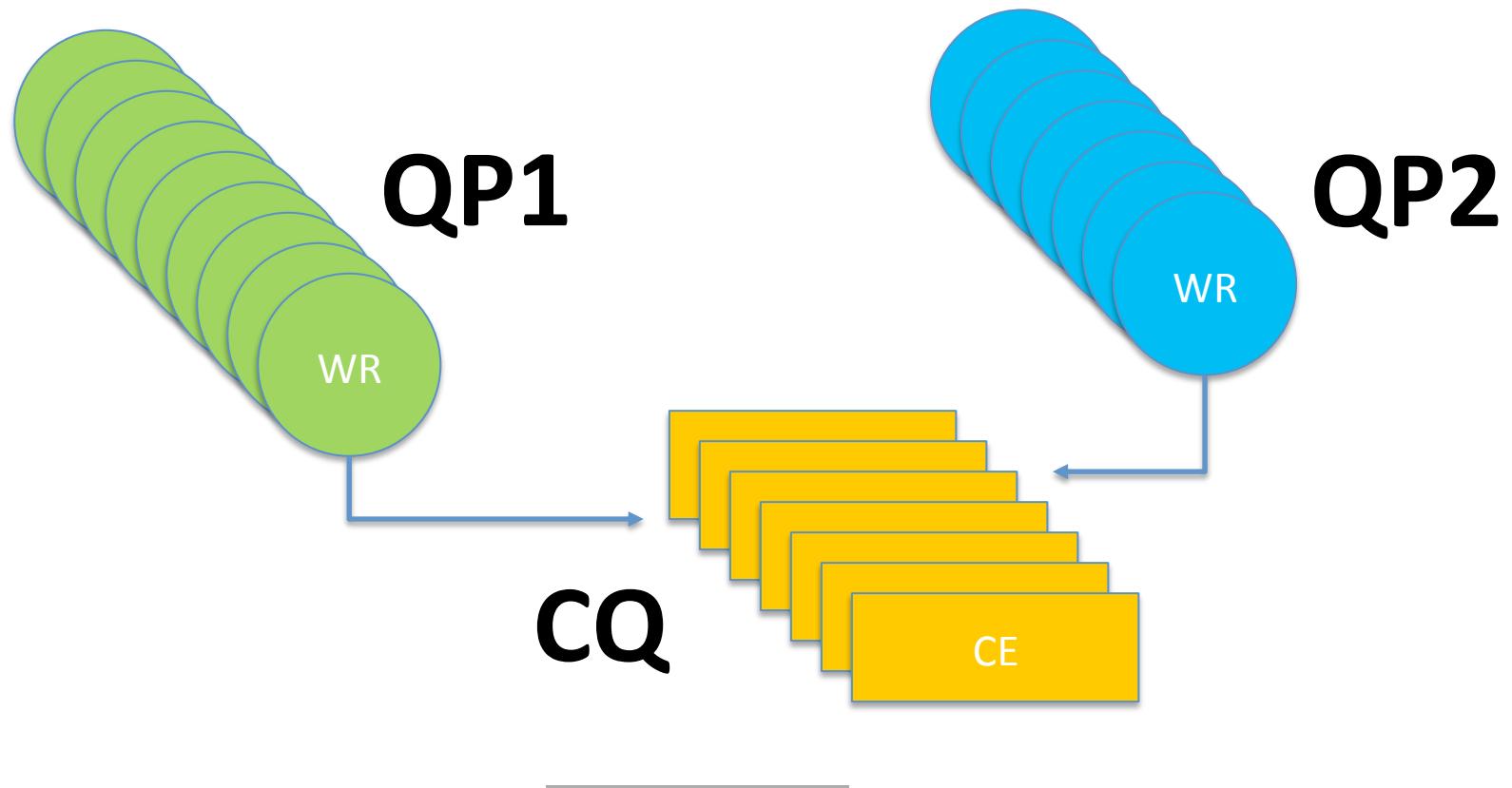
- **Unify completion queue logic from different ULP clients**
- **Simplify completion queue polling and interrupt handling**
  - No need to poll, handle events and maintain logic
- **Resolve the error completions unreliability**
- **Support different polling schemes**
- **Performance optimized**



# DRAINING QP

# PROBLEM

- Unknown when WRs are completed after ceasing posts

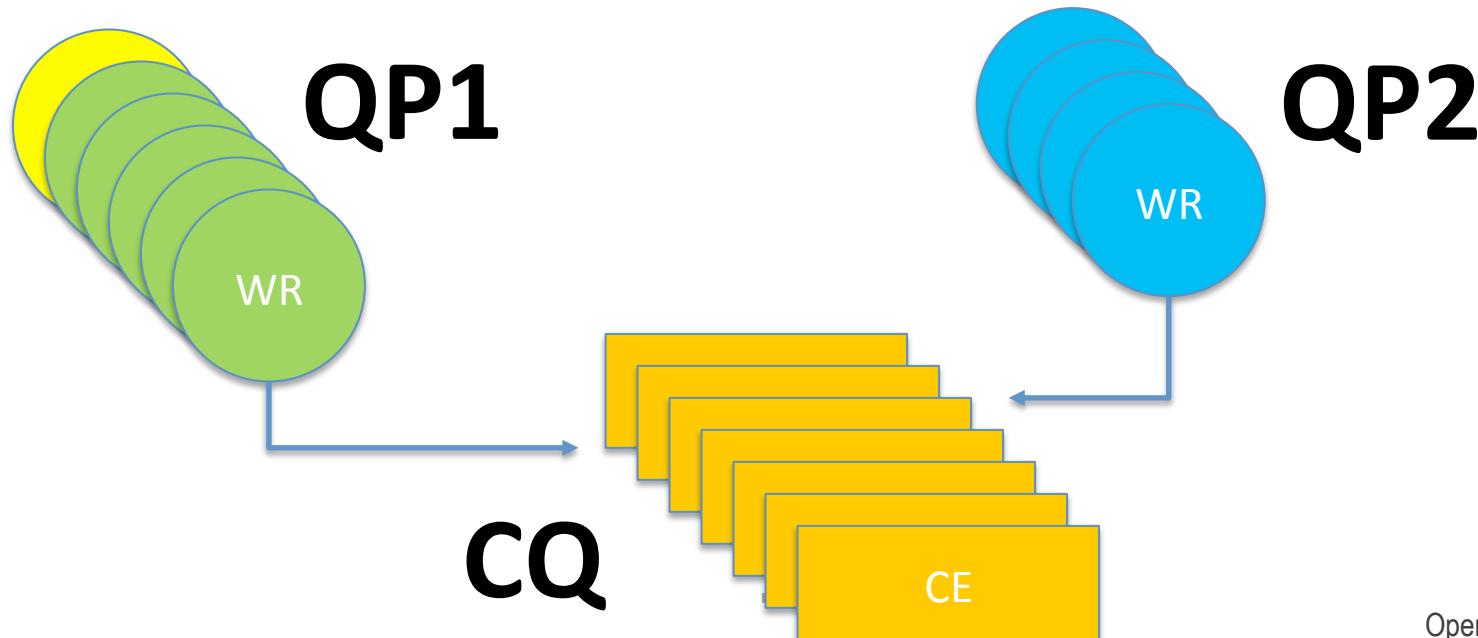


# COMMON METHODS

- 1. Wait until all previously posted WRs complete and then destroy QP - may be indefinite**
- 2. Destroy QP without waiting for completions**
  - need to maintain a shadow state of all in-flight WQEs in order to free related state
- 3. Modify QP to error, and then poll related CQ until it is empty - works only if CQ is associated with only the same QP**

# ROBUST GENERIC DRAINING

1. Cease posting new WRs to QP
2. Change QP state to ERR
3. Post marker WR (nop)
4. Wait until marker WR completes



# DRAIN QP API

OP	API
Drain SQ	void ib_drain_sq(qp)
Drain RQ	void ib_drain_rq(p)
Drain QP	void ib_drain_qp(qp)

- IB/core code to drain SQ/RQ/both in single function call
- Protect from use-after-free error flow
- Synchronized operation in one place



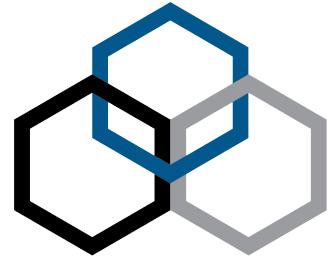
# GENERIC RDMA READ/WRITE (WIP)

# MOTIVATION

- **RDMA READ/WRITE is common operation for storage protocols**
- Every ULP has own implementation
- Lack of support for generic S/G lists
- HCA aware implementations
- Reuse pre-allocated MRs (MR pool)
- Support large number of S/G entries

# REFERENCES

- [IB: new common API for draining queues](#) by Steve Wise
- [Generic RDMA READ/WRITE API](#) by Christoph Hellwig
- [Completion queue abstraction](#) by Christoph Hellwig and Sagi Grimberg
- [New fast registration API](#) by Sagi Grimberg



OPENFABRICS  
ALLIANCE

12<sup>th</sup> ANNUAL WORKSHOP 2016

THANK YOU

Leon Romanovsky

Mellanox Technology

